

OpenFest 2015

Using seccomp to limit the kernel attack surface

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Outline

- 1 Introductions
- 2 Introduction and history
- 3 Seccomp filtering and BPF
- 4 Constructing seccomp filters
- 5 BPF programs
- 6 Further details on seccomp filters
- 7 Applications, tools, and further information

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7 Applications, tools, and further information

Who am I?

- Maintainer of Linux man-pages (since 2004)
 - Documents kernel-user-space + C library APIs
 - ~1000 manual pages
 - <http://www.kernel.org/doc/man-pages/>
- API review, testing, and documentation
 - API design and design review
 - Lots of testing, lots of bug reports, a few kernel patches
- “Day job”: programmer, trainer, writer



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- History of seccomp
- Basics of seccomp operation
- Creating and installing BPF filters (AKA “seccomp2”)
 - Mostly: look at hand-coded BPF filter programs, to gain fundamental understanding of how seccomp works
 - Briefly note some productivity aids for coding BPF programs

Introduction and history

- Mechanism to restrict system calls that a process may make
 - Reduces attack surface of kernel
 - A key component for building application sandboxes
- First version in Linux 2.6.12 (2005)
 - Filtering enabled via `/proc/PID/seccomp`
 - Writing “1” to file places process (irreversibly) in “strict” seccomp mode
 - Need `CONFIG_SECCOMP`



Introduction and history

Initially, just one filtering mode (“strict”)

- Only permitted system calls are *read()*, *write()*, *_exit()*, and *sigreturn()*
 - Note: *open()* not included (must open files before entering strict mode)
 - *sigreturn()* allows for signal handlers
- Other system calls \Rightarrow SIGKILL
- Designed to sandbox compute-bound programs that deal with untrusted byte code
 - Code perhaps exchanged via pre-created pipe or socket



Introduction and history

Linux 2.6.23 (2007):

- `/proc/PID/seccomp` interface replaced by `prctl()` operations
- `prctl(PR_SET_SECCOMP, arg)` modifies caller's seccomp mode
 - `SECCOMP_MODE_STRICT`: limit syscalls as before
- `prctl(PR_GET_SECCOMP)` returns seccomp mode:
 - `0` \Rightarrow process is not in seccomp mode
 - Otherwise?
 - `SIGKILL (!)`
 - `prctl()` is not a permitted system call in "strict" mode
 - Who says kernel developers don't have a sense of humor?



Introduction and history

- Linux 3.5 (2012) adds “filter” mode (AKA “seccomp2”)
 - `prctl(PR_SET_SECCOMP, SECCOMP_MODE_FILTER, ...)`
 - Can control which system calls are permitted,
 - Control based on system call number and argument values
 - Choice is controlled by user-defined filter—a BPF “program”
 - Berkeley Packet Filter (later)
 - Requires `CONFIG_SECCOMP_FILTER`
 - By now used in a range of tools
 - E.g., Chrome browser, OpenSSH, *vsftpd*, systemd, Firefox OS, Docker, LXC



Introduction and history

- Linux 3.8 (2013):
 - The joke is getting old...
 - New `/proc/PID/status` `Seccomp` field exposes process seccomp mode (as a number)

```
0 // SECCOMP_MODE_DISABLED
1 // SECCOMP_MODE_STRICT
2 // SECCOMP_MODE_FILTER
```

- Process can, without fear, read from this file to discover its own seccomp mode
 - But, must have previously obtained a file descriptor...

Linux 3.17 (2014):

- `seccomp()` system call added
 - (Rather than further multiplexing of `prctl()`)
- Provides superset of `prctl(2)` functionality
 - Can synchronize all threads to same filter tree
 - Useful, e.g., if some threads created by start-up code before application has a chance to install filter(s)

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Seccomp filtering and BPF

- Seccomp filtering available since Linux 3.5
- Allows filtering based on system call number and argument (register) values
 - Pointers are **not** dereferenced
- Filters expressed using BPF (Berkeley Packet Filter) syntax
- Filters installed using *seccomp()* or *prctl()*
 - 1 Construct and install BPF filter
 - 2 *exec()* new program or invoke function inside dynamically loaded shared library (plug-in)
- Once installed, **every syscall triggers execution of filter**
 - Installed filters **can't** be removed
 - Filter == declaration that we don't trust subsequently executed code



- BPF originally devised (in 1992) for *tcpdump*
 - Monitoring tool to display packets passing over network
 - <http://www.tcpdump.org/papers/bpf-usenix93.pdf>
- Volume of network traffic is enormous \Rightarrow must filter for packets of interest
- BPF allows **in-kernel selection of packets**
 - Filtering based on fields in packet header
- Filtering in kernel more efficient than filtering in user space
 - Unwanted packets are **discarded early**
 - \Rightarrow Avoids passing **every** packet over kernel-user-space boundary



- BPF defines a **virtual machine** (VM) that can be implemented inside kernel
- VM characteristics:
 - **Simple instruction set**
 - Small set of instructions
 - All instructions are same size
 - Implementation is simple and fast
 - Only **branch-forward** instructions
 - Programs are directed acyclic graphs (DAGs)
 - Easy to verify validity/safety of programs
 - Program completion is guaranteed (DAGs)
 - Simple instruction set \Rightarrow can verify opcodes and arguments
 - Can detect dead code
 - Can verify that program completes via a “return” instruction
 - BPF filter programs are limited to 4096 instructions



Generalizing BPF

- BPF originally designed to work with network packet headers
- Seccomp 2 developers realized BPF could be generalized to solve different problem: filtering of system calls
 - Same basic task: test-and-branch processing based on content of a small set of memory locations
- Further generalization (“extended BPF”) is ongoing
 - Linux 3.18: adding filters to kernel tracepoints
 - Linux 3.19: adding filters to raw sockets
 - In progress (July 2015): filtering of *perf* events



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Basic features of BPF virtual machine

- Accumulator register
- Data area (data to be operated on)
 - In seccomp context: data area describes system call
- Implicit program counter
 - (Recall: all instructions are same size)
- Instructions contained in structure of this form:

```
struct sock_filter { /* Filter block */
    __u16 code;      /* Filter code (opcode)*/
    __u8  jt;        /* Jump true */
    __u8  jf;        /* Jump false */
    __u32 k;         /* Generic multiuse field */
};
```

- See `<linux/filter.h>` and `<linux/bpf_common.h>`



Instruction set includes:

- Load instructions
- Store instructions
- Jump instructions
- Arithmetic/logic instructions
 - ADD, SUB, MUL, DIV, MOD, NEG
 - OR, AND, XOR, LSH, RSH
- Return instructions
 - Terminate filter processing
 - Report a status telling kernel what to do with syscall

BPF jump instructions

- Conditional and unconditional jump instructions provided
- Conditional jump instructions consist of
 - **Opcode** specifying condition to be tested
 - **Value** to test against
 - **Two** jump targets
 - **jt**: target if condition is true
 - **jf**: target if condition is false
- Conditional jump instructions:
 - **JEQ**: jump if equal
 - **JGT**: jump if greater
 - **JGE**: jump if greater or equal
 - **JSET**: bit-wise AND + jump if nonzero result
 - **jf** target \Rightarrow no need for **JNE**, **JLT**, **JLE**, and **JCLEAR**



BPF jump instructions

- Targets are expressed as relative offsets in instruction list
 - 0 == no jump (execute next instruction)
 - `jt` and `jf` are 8 bits \Rightarrow 255 maximum offset for conditional jumps
- Unconditional `JA` (“jump always”) uses `k` as offset, allowing much larger jumps

Seccomp BPF data area

- Seccomp provides data describing syscall to filter program
 - Buffer is **read-only**
- Format (expressed as C `struct`):

```
struct seccomp_data {
    int    nr;                /* System call number */
    __u32  arch;             /* AUDIT_ARCH_* value */
    __u64  instruction_pointer; /* CPU IP */
    __u64  args[6];         /* System call arguments */
};
```



Seccomp BPF data area

```
struct seccomp_data {
    int    nr;                /* System call number */
    __u32  arch;             /* AUDIT_ARCH_* value */
    __u64  instruction_pointer; /* CPU IP */
    __u64  args[6];         /* System call arguments */
};
```

- *nr*: system call number (architecture-dependent)
- *arch*: identifies architecture
 - Constants defined in `<linux/audit.h>`
 - `AUDIT_ARCH_X86_64`, `AUDIT_ARCH_I386`,
`AUDIT_ARCH_ARM`, etc.
- *instruction_pointer*: CPU instruction pointer
- *args*: system call arguments
 - System calls have maximum of six arguments
 - Number of elements used depends on system call



Building BPF instructions

- Obviously, one can code BPF instructions numerically by hand
- But, header files define symbolic constants and convenience macros (`BPF_STMT()`, `BPF_JUMP()`) to ease the task

```
#define BPF_STMT(code, k) \  
    { (unsigned short)(code), 0, 0, k }  
#define BPF_JUMP(code, k, jt, jf) \  
    { (unsigned short)(code), jt, jf, k }
```

- (Macros just plug values together to form structure)



Building BPF instructions: examples

- Load architecture number into accumulator

```
BPF_STMT(BPF_LD | BPF_W | BPF_ABS ,  
         (offsetof(struct seccomp_data, arch)))
```

- Opcode here is constructed by ORing three values together:
 - **BPF_LD**: load
 - **BPF_W**: operand size is a word
 - **BPF_ABS**: address mode specifying that source of load is data area (containing system call data)
 - See `<linux/bpf_common.h>` for definitions of opcode constants
- `offsetof()` generates offset of desired field in data area



Building BPF instructions: examples

- Test value in accumulator

```
BPF_JUMP(BPF_JMP | BPF_JEQ | BPF_K,  
         AUDIT_ARCH_X86_64, 1, 0)
```

- **BPF_JMP | BPF_JEQ**: jump with test on equality
- **BPF_K**: value to test against is in generic multiuse field (**k**)
- **k** contains value **AUDIT_ARCH_X86_64**
- **jt** value is 1, meaning skip one instruction if test is true
- **jf** value is 0, meaning skip zero instructions if test is false
 - I.e., continue execution at following instruction
- Return value that causes kernel to kill process with **SIGSYS**

```
BPF_STMT(BPF_RET | BPF_K, SECCOMP_RET_KILL)
```



Checking the architecture

- Checking architecture value should be first step in any BPF program
- Architecture may support multiple system call conventions
 - E.g. x86 hardware supports x86-64 and i386
 - System call numbers may differ or overlap

Filter return value

- Once a filter is installed, each system call is tested against filter
- Seccomp filter must return a value to kernel indicating whether system call is permitted
 - Otherwise `EINVAL` when attempting to install filter
- Return value is 32 bits, in two parts:
 - Most significant 16 bits (`SECCOMP_RET_ACTION` mask) specify an action to kernel
 - Least significant 16 bits (`SECCOMP_RET_DATA` mask) specify “data” for return value



Filter return action

Filter return action component is one of

- `SECCOMP_RET_ALLOW`: system call is executed
- `SECCOMP_RET_KILL`: process is immediately terminated
 - Terminated *as though* process had been killed with `SIGSYS`
- `SECCOMP_RET_ERRNO`: return an error from system call
 - System call is not executed
 - Value in `SECCOMP_RET_DATA` is returned in `errno`
- `SECCOMP_RET_TRACE`: attempt to notify `ptrace()` tracer
 - Gives tracing process a chance to assume control
 - See `seccomp(2)`
- `SECCOMP_RET_TRAP`: process is sent `SIGSYS` signal
 - Can catch this signal; see `seccomp(2)` for more details



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Installing a BPF program

- A process installs a filter for itself using one of:
 - `seccomp(SECCOMP_SET_MODE_FILTER, flags, &fprog)`
 - Only since Linux 3.17
 - `prctl(PR_SET_SECCOMP, SECCOMP_MODE_FILTER, &fprog)`
- *&fprog* is a pointer to a BPF program:

```
struct sock_fprog {
    unsigned short len; /* Number of instructions */
    struct sock_filter *filter;
                        /* Pointer to program
                           (array of instructions) */
};
```



Installing a BPF program

To install a filter, one of the following must be true:

- Caller is privileged (`CAP_SYS_ADMIN`)
- Caller has to set the `no_new_privs` process attribute:

```
prctl(PR_SET_NO_NEW_PRIVS, 1, 0, 0, 0);
```

- Causes set-UID/set-GID bit / file capabilities to be ignored on subsequent `execve()` calls
 - Once set, `no_new_privs` can't be unset
- Prevents possibility of attacker starting privileged program and manipulating it to misbehave using a seccomp filter
- `! no_new_privs && ! CAP_SYS_ADMIN ⇒ seccomp()` fails with `EACCES`



Example: seccomp/seccomp_deny_open.c

```
1 int main(int argc, char **argv) {
2     prctl(PR_SET_NO_NEW_PRIVS, 1, 0, 0, 0);
3
4     install_filter();
5
6     open("/tmp/a", O_RDONLY);
7
8     printf("We shouldn't see this message\n");
9     exit(EXIT_SUCCESS);
10 }
```

Program installs a filter that prevents *open()* being called, and then calls *open()*

- Set `no_new_privs` bit
- Install seccomp filter
- Call *open()*



Example: seccomp/seccomp_deny_open.c

```
1 static void install_filter(void) {
2     struct sock_filter filter[] = {
3         BPF_STMT(BPF_LD | BPF_W | BPF_ABS,
4                 (offsetof(struct seccomp_data, arch))),
5         BPF_JUMP(BPF_JMP | BPF_JEQ | BPF_K,
6                 AUDIT_ARCH_X86_64, 1, 0),
7         BPF_STMT(BPF_RET | BPF_K, SECCOMP_RET_KILL),
8         ...
9     }
```

- Define and initialize array (of structs) containing BPF filter program
- Load architecture into accumulator
- Test if architecture value matches `AUDIT_ARCH_X86_64`
 - True: jump forward one instruction (i.e., skip next instruction)
 - False: skip no instructions
- Kill process on architecture mismatch



Example: seccomp/seccomp_deny_open.c

```
1     BPF_STMT(BPF_LD | BPF_W | BPF_ABS,  
2             (offsetof(struct seccomp_data, nr))),  
3  
4     BPF_JUMP(BPF_JMP | BPF_JEQ | BPF_K, __NR_open,  
5             1, 0),  
6     BPF_STMT(BPF_RET | BPF_K, SECCOMP_RET_ALLOW),  
7  
8     BPF_STMT(BPF_RET | BPF_K, SECCOMP_RET_KILL)  
9 };
```

- Remainder of filter program
- Load system call number into accumulator
- Test if system call number matches `__NR_open`
 - True: advance one instruction \Rightarrow kill process
 - False: advance 0 instructions \Rightarrow allow system call



Example: seccomp/seccomp_deny_open.c

```
1  struct sock_fprog prog = {
2      .len = (unsigned short) (sizeof(filter) /
3                              sizeof(filter[0])),
4      .filter = filter,
5  };
6
7  seccomp(SECCOMP_SET_MODE_FILTER, 0, &prog);
8 }
```

- Construct argument for *seccomp()*
- Install filter

Example: seccomp/seccomp_deny_open.c

Upon running the program, we see:

```
$ ./seccomp_deny_open
Bad system call      # Message printed by shell
$ echo $?           # Display exit status of last command
159
```

- “Bad system call” indicates process was killed by **SIGSYS**
- Exit status of 159 ($== 128 + 31$) also indicates termination as though killed by **SIGSYS**
 - Exit status of process killed by signal is $128 + \textit{signum}$
 - **SIGSYS** is signal number 31 on this architecture

Example: `seccomp/seccomp_control_open.c`

- A more sophisticated example
- Filter based on *flags* argument of *open()*
 - `O_CREAT` specified \Rightarrow kill process
 - `O_WRONLY` or `O_RDWR` specified \Rightarrow cause *open()* to fail with `ENOTSUP` error

Example: seccomp/seccomp_control_open.c

```
struct sock_filter filter[] = {
    BPF_STMT(BPF_LD | BPF_W | BPF_ABS,
             (offsetof(struct seccomp_data, arch))),
    BPF_JUMP(BPF_JMP | BPF_JEQ | BPF_K,
            AUDIT_ARCH_X86_64, 1, 0),
    BPF_STMT(BPF_RET | BPF_K, SECCOMP_RET_KILL),

    BPF_STMT(BPF_LD | BPF_W | BPF_ABS,
             (offsetof(struct seccomp_data, nr))),
    BPF_JUMP(BPF_JMP | BPF_JEQ | BPF_K, __NR_open, 1, 0),
    BPF_STMT(BPF_RET | BPF_K, SECCOMP_RET_ALLOW),
```

- Load architecture and test for expected value
- Load system call number
- Test if system call number is `__NR_open`
 - True: skip next instruction
 - False: skip 0 instructions \Rightarrow permit all other syscalls



Example: seccomp/seccomp_control_open.c

```
BPF_STMT(BPF_LD | BPF_W | BPF_ABS ,
          (offsetof(struct seccomp_data, args[1]))),

BPF_JUMP(BPF_JMP | BPF_JSET | BPF_K, O_CREAT, 0, 1),
BPF_STMT(BPF_RET | BPF_K, SECCOMP_RET_KILL),
```

- Load second argument of `open()` (`flags`)
- Test if `O_CREAT` bit is set in `flags`
 - True: skip 0 instructions \Rightarrow kill process
 - False: skip 1 instruction

Example: seccomp/seccomp_control_open.c

```
BPF_JUMP(BPF_JMP | BPF_JSET | BPF_K,  
         O_WRONLY | O_RDWR, 0, 1),  
BPF_STMT(BPF_RET | BPF_K,  
         SECCOMP_RET_ERRNO |  
         (ENOTSUP & SECCOMP_RET_DATA)),  
BPF_STMT(BPF_RET | BPF_K, SECCOMP_RET_ALLOW)  
};
```

- Test if `O_WRONLY` or `O_RDWR` are set in *flags*
 - True: cause `open()` to fail with `ENOTSUP` error in *errno*
 - False: allow `open()` to proceed



Example: `seccomp/seccomp_control_open.c`

```
int main(int argc, char **argv) {
    prctl(PR_SET_NO_NEW_PRIVS, 1, 0, 0, 0);
    install_filter();

    if (open("/tmp/a", O_RDONLY) == -1)
        perror("open1");
    if (open("/tmp/a", O_WRONLY) == -1)
        perror("open2");
    if (open("/tmp/a", O_RDWR) == -1)
        perror("open3");
    if (open("/tmp/a", O_CREAT | O_RDWR, 0600) == -1)
        perror("open4");

    exit(EXIT_SUCCESS);
}
```

- Test `open()` calls with various flags



Example: seccomp/seccomp_control_open.c

```
$ ./seccomp_control_open
open2: Operation not supported
open3: Operation not supported
Bad system call
$ echo $?
159
```

- First *open()* succeeded
- Second and third *open()* calls failed
 - Kernel produced `ENOTSUP` error for call
- Fourth *open()* call caused process to be killed

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Installing multiple filters

- If existing filters permit `prctl()` or `seccomp()`, further filters can be installed
- **All** filters are always executed, in **reverse order** of registration
- Each filter yields a return value
- Value returned to kernel is first seen action of highest priority (along with accompanying data)
 - `SECCOMP_RET_KILL` (highest priority)
 - `SECCOMP_RET_TRAP`
 - `SECCOMP_RET_ERRNO`
 - `SECCOMP_RET_TRACE`
 - `SECCOMP_RET_ALLOW` (lowest priority)



fork() and *execve()* semantics

- If seccomp filters permit *fork()* or *clone()*, then child inherits parents filters
- If seccomp filters permit *execve()*, then filters are preserved across *execve()*



Cost of filtering, construction of filters

- Installed BPF filter(s) are executed for every system call
 - \Rightarrow there's a performance cost
- Example on x86-64:
 - Use our “deny open” seccomp filter
 - Requires 6 BPF instructions / permitted syscall
 - Call `getppid()` repeatedly (one of cheapest syscalls)
 - +25% execution time (with JIT compiler disabled)
 - (Looks relatively high because `getppid()` is a cheap syscall)
- Obviously, order of filtering rules can affect performance
 - Construct filters so that most common cases yield shortest execution paths
 - If handling many different system calls, binary chop techniques can give $O(\log N)$ performance



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Possible applications:

- Building sandboxed environments
 - Whitelisting usually safer than blacklisting
 - Default treatment: block all system calls
 - Then allow only a limited set of syscall / argument combinations
 - Various examples mentioned earlier
- Failure-mode testing
 - Place application in environment where unusual / unexpected failures occur
 - Blacklist certain syscalls / argument combinations to generate failures

- High-level API for kernel creating seccomp filters
 - <https://github.com/seccomp/libseccomp>
 - Initial release: 2012
- Simplifies various aspects of building filters
 - Eliminates tedious/error-prone tasks such as changing branch instruction counts when instructions are inserted
 - Abstract architecture-dependent details out of filter creation
 - Can output generated code in binary (for seccomp filtering) or human-readable form (“pseudofilter code”)
 - Don’t have full control of generated code, but can give hints about which system calls to prioritize in generated code
- <http://lwn.net/Articles/494252/>
- Fully documented with man pages that contain examples (!)



- *bpfc* (BPF compiler)
 - Compiles assembler-like BPF programs to byte code
 - Part of *netsniff-ng* project (<http://netsniff-ng.org/>)
- LLVM has an eBPF back end (merged Jan 2015)
 - eBPF support for seccomp is planned
 - Compiles subset of C to BPF
 - C dialect; does not provide: loops, global variables, FP numbers, vararg functions, passing structs as args...
 - Examples in kernel source: `samples/bpf/*_kern.c`
 - GCC patches exist, but not (yet?) merged upstream

- In-kernel JIT (just-in-time) compiler
 - Compiles BPF binary to native machine code at load time
 - Execution speed up of 2x to 3x (or better, in some cases)
 - Disabled by default; enable by writing “1” to `/proc/sys/net/core/bpf_jit_enable`
 - See [bpf\(2\)](#) man page

- Kernel source files:
`Documentation/prctl/seccomp_filter.txt`,
`Documentation/networking/filter.txt`
- <http://outflux.net/teach-seccomp/>
 - Shows handy trick for discovering which of an application's system calls don't pass filtering
- [seccomp\(2\)](#) man page
- "Seccomp sandboxes and memcached example"
 - blog.viraptor.info/post/seccomp-sandboxes-and-memcached-example-part-1
 - blog.viraptor.info/post/seccomp-sandboxes-and-memcached-example-part-2
- <https://lwn.net/Articles/656307/>
 - Write-up of a version of this presentation...



Thanks!

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